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# Managing a Fragmented XML Data Cube with Oracle and Timesten

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#### **Outline**

- Context and Motivation
- Timesten in-Memory Database
- XML Data Cube
- XML Cube Management Configurations
- Implementation and Testing
- Related Work
- Conclusion

#### Context and motivation

- Performance optimization of data warehouses (DW)
- Focus on a special type of DWs: XML warehouses
- Warehousing and analyzing complex data
  - Multidimensional model: conceptual, logical, physical level
  - Performance issues
  - Vertical fragmentation approach proposed
- Crossing two techniques
  - Vertical Fragmentation
  - Caching

# Objectives and Contribution

#### Objective:

- Analyze the impact of vertical fragmentation on caching and vice versa
- A better cache management (data organization-aware)
- Leverage the vertical fragmentation

#### Contributions

- A set of configurations to manage a fragmented XML cube
- A comparision between the configurations

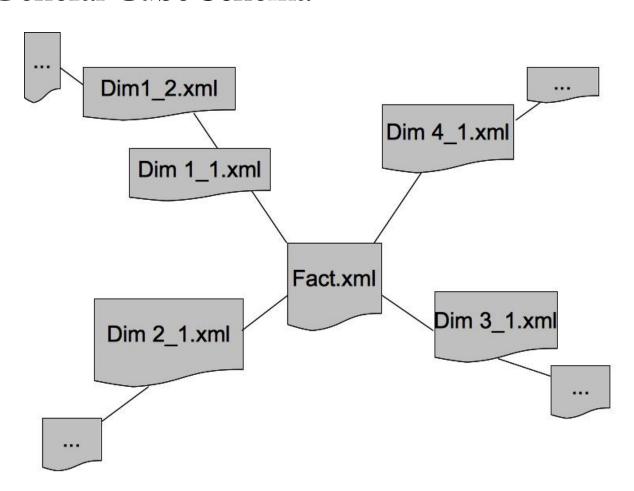
# TimesTen In-Memory database

- Oracle's In-memory database solution
- Different uses of TimesTen

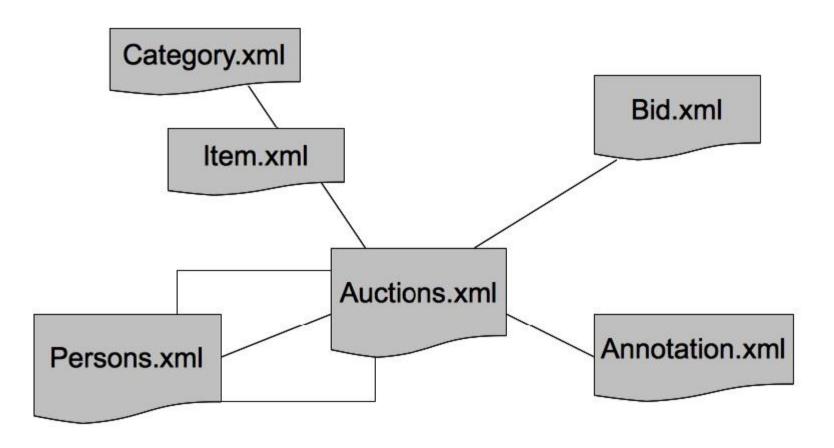
- As a database cache for a disk resident database
  - Read-only transactions
  - Read-Write transactions
- As a full-featured relational database
  - Persistence
  - Recovery
  - **...**

#### XML cubeModel

General Cube Schema

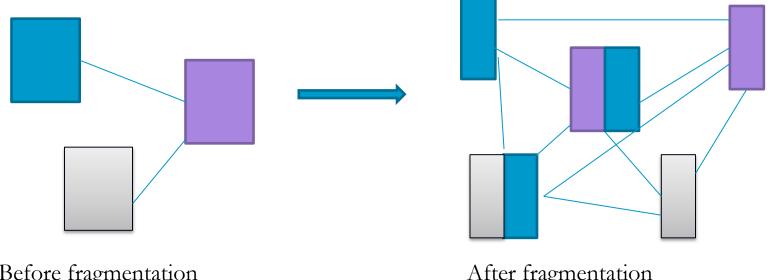


■ Instantiation: Auction cube



- Unfragmented XML cube
  - Basic configuration
  - Fact and each dimension member = one XML document
  - Formally:
    - UXCube =  $\{D_i, i=1, ...\}$  set of XML documents
    - Di =  $\{P_j^i, i=1, ..., j=1,...\}$  set of XML properties accessed by XPath

- Fragmented XML cube
  - Vertical fragmentation approach for XML Cubes (Dawak'11)
  - Each document of the Cube split into fragments
    - Homogeneous fragments: properties ∈ same original document
    - Heteregenous fragments: properties



Before fragmentation

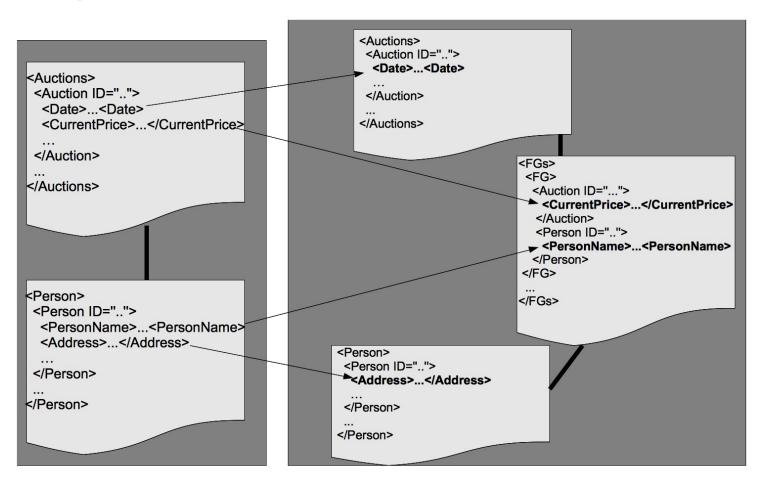
After fragmentation

- Fragmented XML cube
  - Frequent Fragment: derived from a frequent property set (Association rules)
  - InFrequent Fragment: properties ¬∈frequent property set

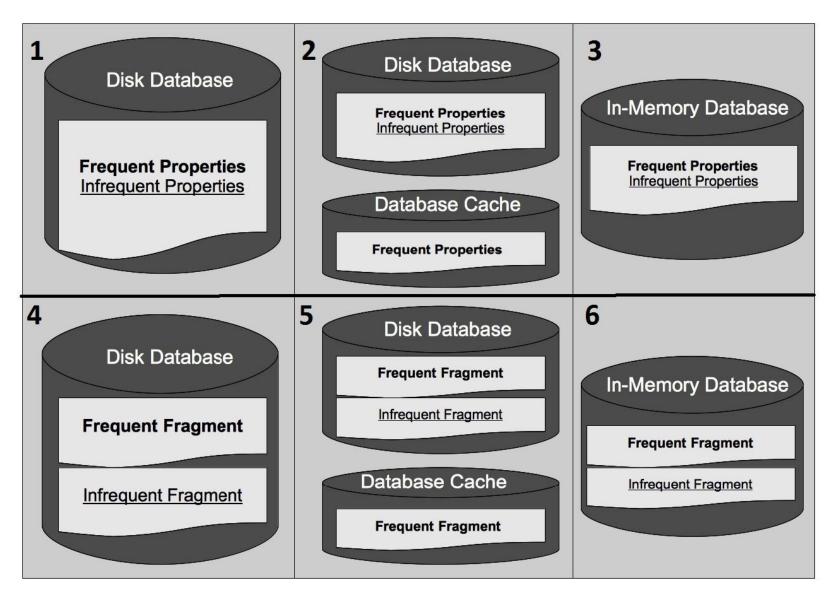
Formally

$$FXCube = \{FF_m, m=1, ...\} \cup \{IF_n, n=1, ...\}$$

Fragmented XML cube: example

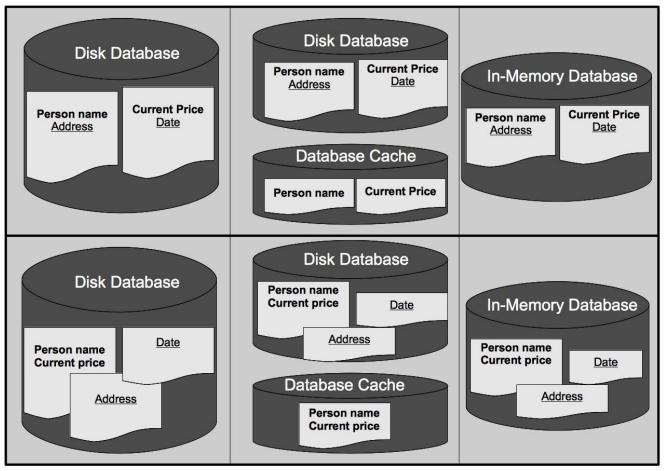


### XML Cube Management Configurations



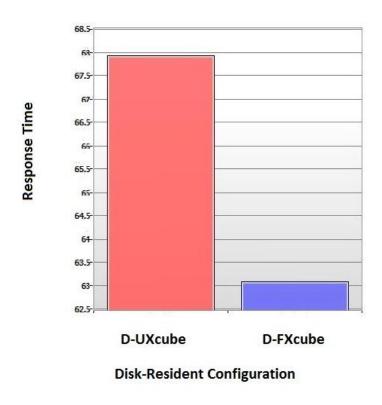
### XML Cube Management Configurations

Instantiation: Auction cube configurations

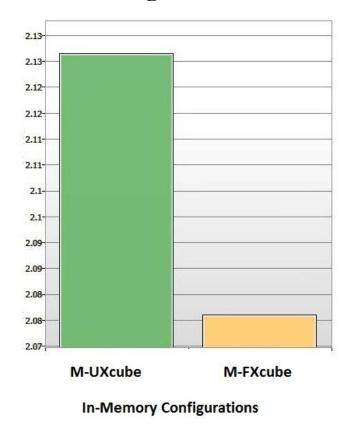


- Disk Resident Database: Oracle 11g Rel. 2
- Database cache and in-Memory database: Oracle TimesTen 11.2.1
- Data Set:
  - XML Cube of auctions: 6 XML document types
  - Fragmented Cube: 28 XML fragment types
- Query load: 100 analytical queries targeting different aggregation levels of UXCube
- Queries rewritten against FXCube

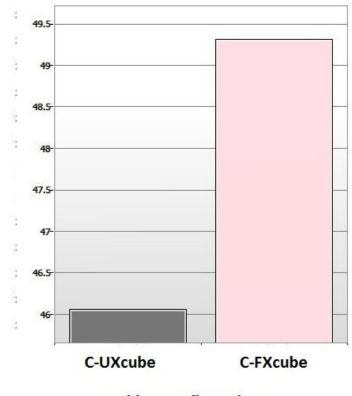
- First measure: average query response times
  - Unfragmented Vs Fragmented XML Cube



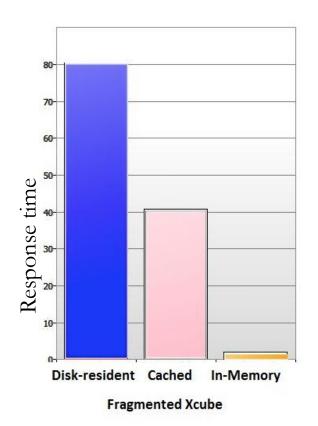
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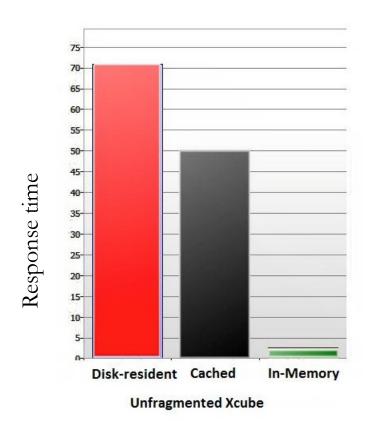
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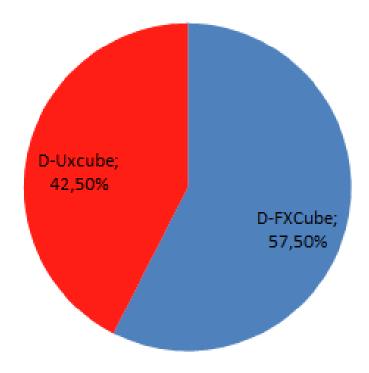
- First measure: average query response times
  - Disk-resident Vs Cached Vs in-Memory XML Cube



- First measure: average query response times
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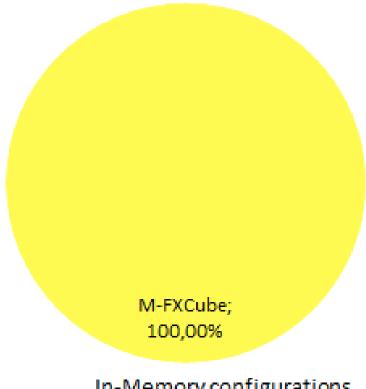


- Second measure: percentage of efficient queries
  - Unfragmented Vs Fragmented XML Cube



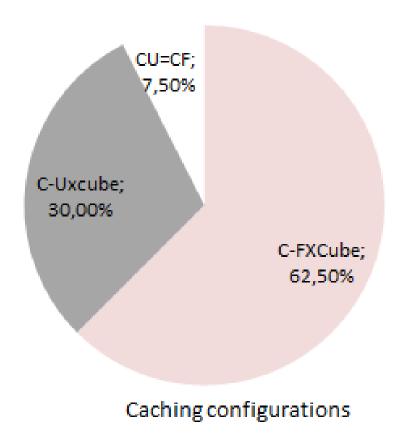
Disk resident configurations

- Second measure: percentage of efficient queries
  - Unfragmented Vs Fragmented XML Cube

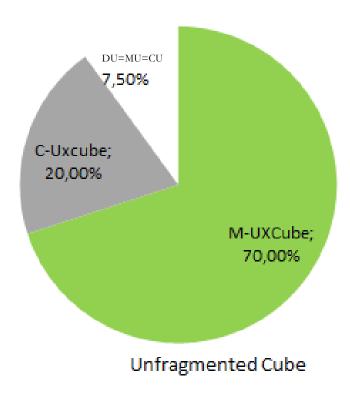


In-Memory configurations

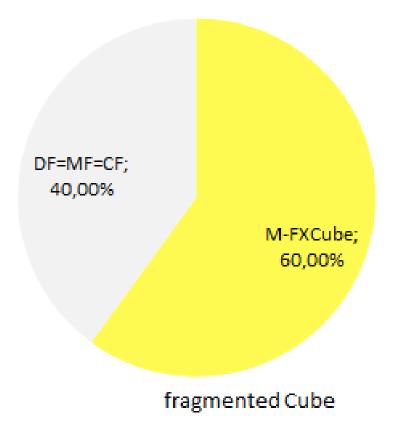
- Second measure: percentage of efficient queries
  - Unfragmented Vs Fragmented XML Cube



- Second measure: percentage of efficient queries
  - Disk-resident Vs Cached Vs in-Memory XML Cube



- Second measure: percentage of efficient queries
  - Disk-resident Vs Cached Vs in-Memory XML Cube



# Related Work (1/2)

Category of work	Examples
Database & Web	<ul> <li>Altinel et al (2003): Static and dynamic caching</li> <li>Manegold et al (2000): Optimizing main memory access</li> <li>Dar et al. (1996): Semantic caching</li> <li>Huang and Hsu (2008): Web document caching</li> </ul>
Data Warehouses	<ul> <li>Andrade et al. (2007): Optimizing multiple data analysis queries</li> <li>Deshpande et al. (1998) Cache small regions of a multidimensional space</li> <li>Lehner et al. (2000): Dynamic caching for multidimensional data</li> <li>Scheuermann et al. (1996): Caching small sets of query results</li> <li>Muto &amp; Kitsuregawa (1998): Main memory for compressed cube management</li> <li>Ross &amp; Zaman (2000): Cache data cube subset materialization</li> </ul>
XML	<ul> <li>Yang et al. (2003): Cache frequent XML patterns</li> <li>Mandhani &amp; Suciu (2005): Semantic cache of materialized Xpath queries</li> <li>Obermeier and Bottcher (2008): XML splitting over mobile devices</li> </ul>

# Related Work (2/2)

Discussion

 Our work meets the same motivation of Obermeier and Bottcher (2008) but applied to XML cubes

 Combination of Vertical fragmentation, main memory data management and caching not tackled before

# Conclusion and Perspectives

#### Conclusion

- Crossed two optimization techniques of data warehouses:
   caching and vertical fragmentation
- Benefits of fragmentation when the cube is managed in main memory
- In-memory enhances both fragmented and unfragmented cube
- Main memory increases the % of efficient queries

# Conclusion and Perspectives

- Pespectives
  - Implement our proposals on an ad hoc network
  - Combine horizontal and vertical fragmentation with cache and in-memory management